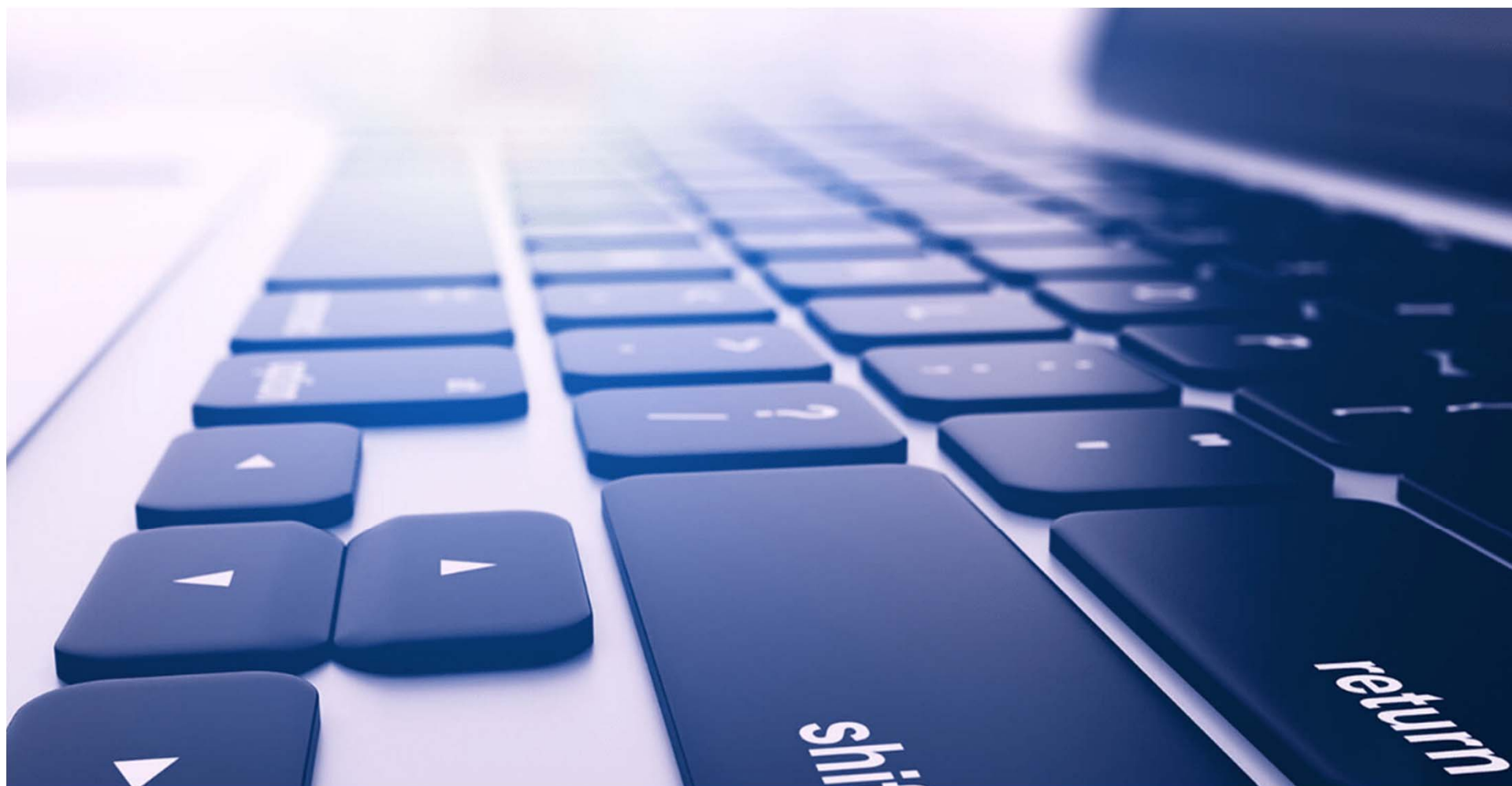



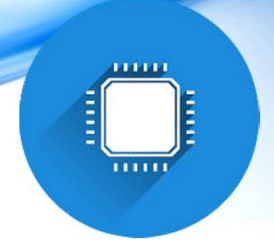


معماری کامپیوتر

جلسه یازدهم: سلسله مراتب حافظه  
فصل هفت کتاب پترسن

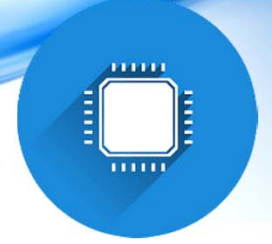


- 
- انواع حافظه
  - سلسله مراتب حافظه
  - مفهوم محلی بودن
    - زمانی
    - مکانی
  - نگاشت مستقیم
  - مثال های متنوع
  - نگاشت مجموعه انجمنی (اشتراکی)

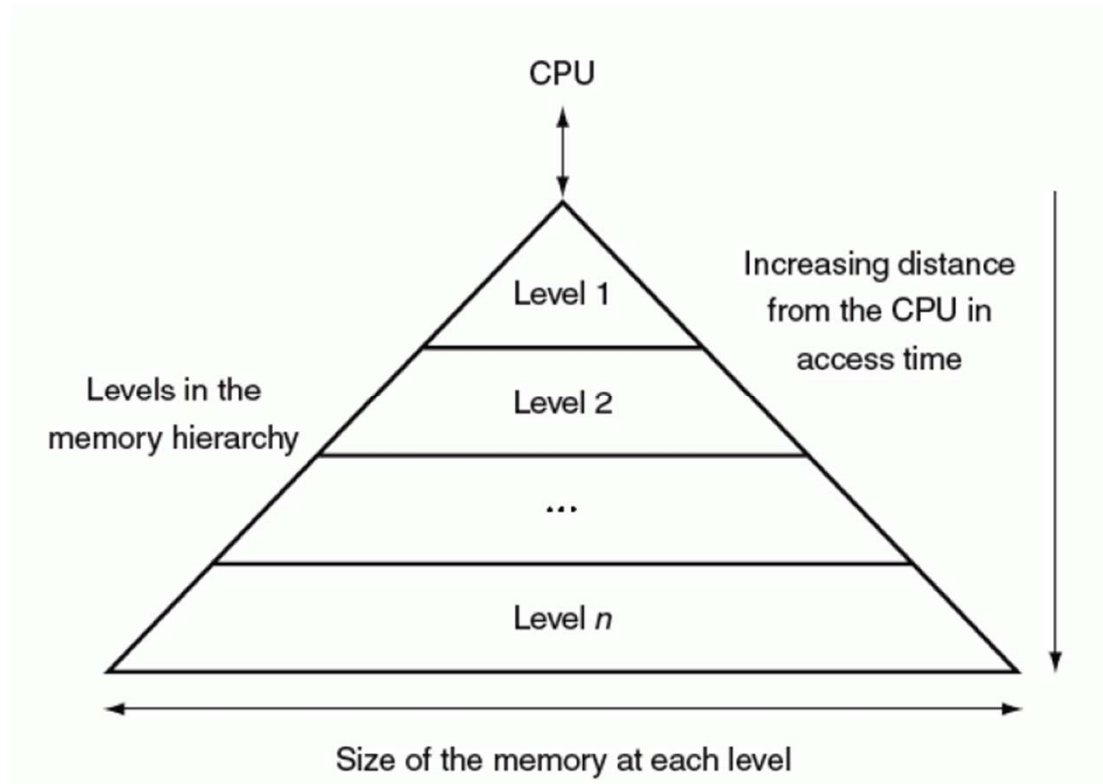


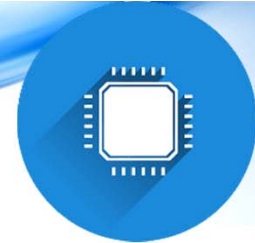
## انواع حافظه

Speed	CPU	Size	Cost (\$/bit)	Current Technology
Fastest	Memory	Smallest	Highest	SRAM
	Memory			DRAM
Slowest	Memory	Biggest	Lowest	Magnetic Disk



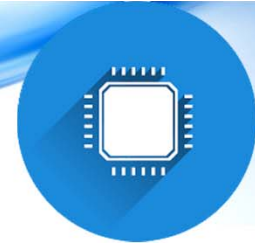
## سلسله مراتب حافظه



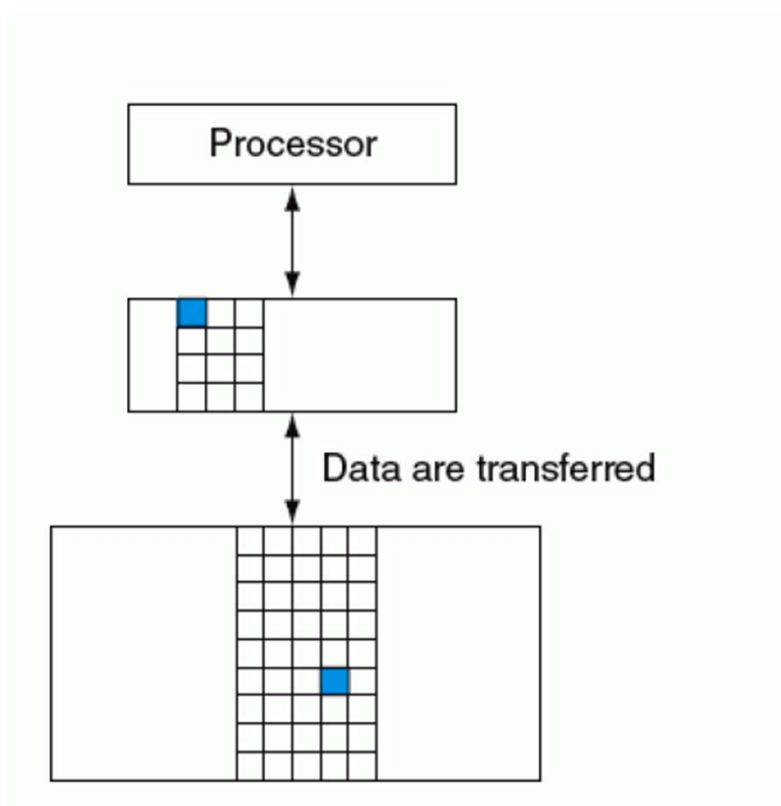


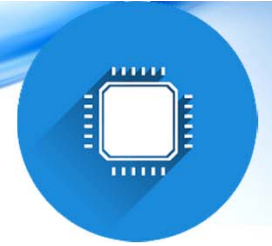
## محلی بودن ارجاعات زمانی و مکانی

- **Temporal locality** (locality in time): If an item is referenced, it will tend to be referenced again soon. If you recently brought a book to your desk to look at, you will probably need to look at it again soon.
- **Spatial locality** (locality in space): If an item is referenced, items whose addresses are close by will tend to be referenced soon. For example, when

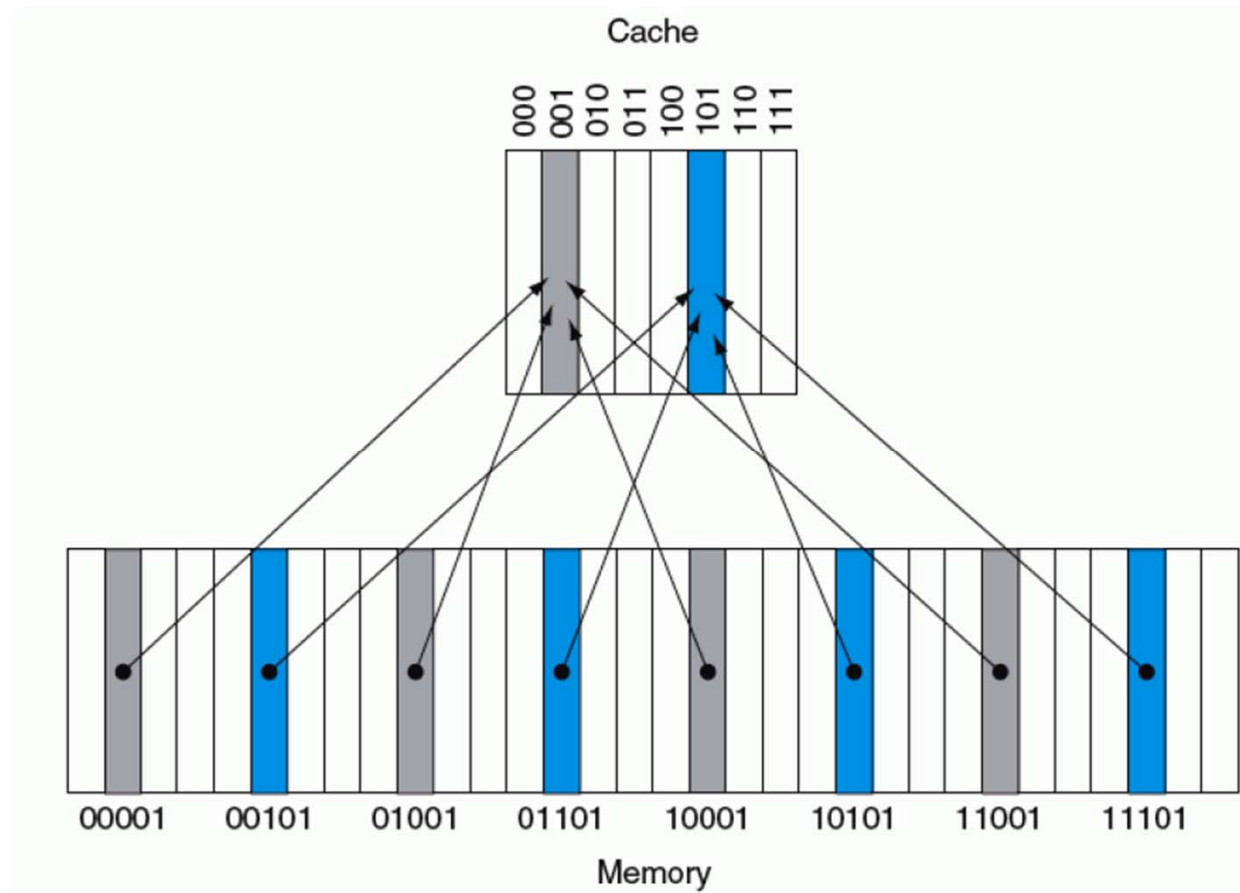


## مفهوم حافظه نهان

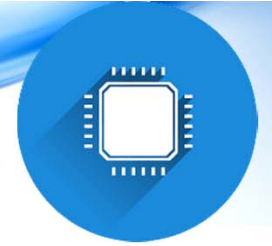




## نگاشت مستقیم



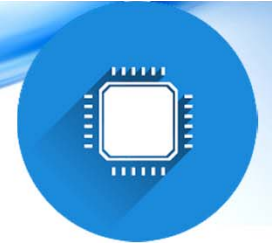




## وضعیت اولیه

Index	V	Tag	Data
000	N		
001	N		
010	N		
011	N		
100	N		
101	N		
110	N		
111	N		

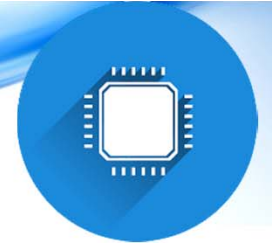
a. The initial state of the cache after power-on



بعد از اولین درخواست 10110

Index	V	Tag	Data
000	N		
001	N		
010	N		
011	N		
100	N		
101	N		
110	Y	$10_{two}$	Memory( $10110_{two}$ )
111	N		

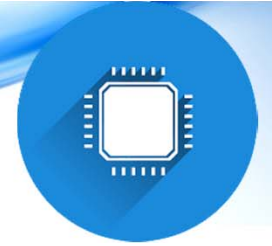
b. After handling a miss of address ( $10110_{two}$ )



بعد از درخواست 11010

Index	V	Tag	Data
000	N		
001	N		
010	Y	$11_{two}$	Memory ( $11010_{two}$ )
011	N		
100	N		
101	N		
110	Y	$10_{two}$	Memory ( $10110_{two}$ )
111	N		

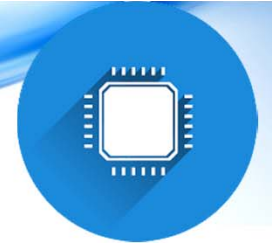
c. After handling a miss of address ( $11010_{two}$ )



بعد از درخواست 10000

Index	V	Tag	Data
000	Y	$10_{two}$	Memory ( $10000_{two}$ )
001	N		
010	Y	$11_{two}$	Memory ( $11010_{two}$ )
011	N		
100	N		
101	N		
110	Y	$10_{two}$	Memory ( $10110_{two}$ )
111	N		

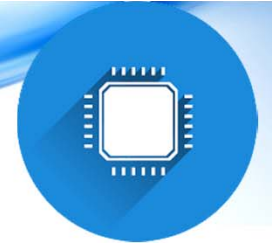
d. After handling a miss of address ( $10000_{two}$ )



بعد از درخواست 00011

Index	V	Tag	Data
000	Y	$10_{two}$	Memory ( $10000_{two}$ )
001	N		
010	Y	$11_{two}$	Memory ( $11010_{two}$ )
011	Y	$00_{two}$	Memory ( $00011_{two}$ )
100	N		
101	N		
110	Y	$10_{two}$	Memory ( $10110_{two}$ )
111	N		

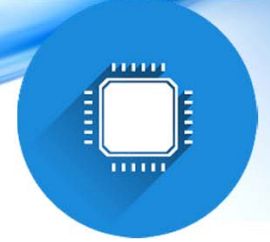
e. After handling a miss of address ( $00011_{two}$ )



بعد از درخواست 10010

Index	V	Tag	Data
000	Y	$10_{\text{two}}$	Memory ( $10000_{\text{two}}$ )
001	N		
010	Y	$10_{\text{two}}$	Memory ( $10010_{\text{two}}$ )
011	Y	$00_{\text{two}}$	Memory ( $00011_{\text{two}}$ )
100	N		
101	N		
110	Y	$10_{\text{two}}$	Memory ( $10110_{\text{two}}$ )
111	N		

f. After handling a miss of address ( $10010_{\text{two}}$ )



## مثال از نگاهت مستقیم

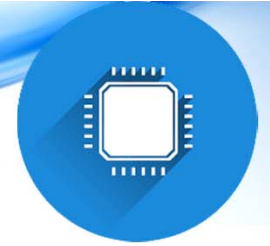
### Bits in a Cache

How many total bits are required for a direct-mapped cache with 16 KB of data and 4-word blocks, assuming a 32-bit address? هر کلمه ۴ بایت است

We know that 16 KB is 4K words, which is  $2^{12}$  words, and, with a block size of 4 words ( $2^2$ ),  $2^{10}$  blocks. Each block has  $4 \times 32$  or 128 bits of data plus a tag, which is  $32 - 10 - 2 - 2$  bits, plus a valid bit. Thus, the total cache size is

$$2^{10} \times (128 + (32 - 10 - 2 - 2) + 1) = 2^{10} \times 147 = 147 \text{ Kbits}$$

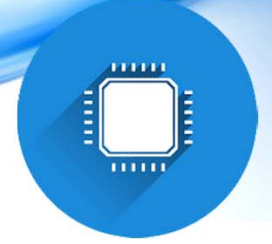
or 18.4 KB for a 16 KB cache. For this cache, the total number of bits in the cache is about 1.15 times as many as needed just for the storage of the data.



## مثال از نگاشت مستقیم

- *How many total bits are required for a direct-mapped cache with 128 KB of data and 1-word block size, assuming a 32-bit address?* هر کلمه ۴ بایت است
- Cache data = 128 KB =  $2^{17}$  bytes =  $2^{15}$  words =  $2^{15}$  blocks
- Cache entry size = block data bits + tag bits + valid bit  
=  $32 + (32 - 15 - 2) + 1 = 48$  bits
- Therefore, cache size =  $2^{15} \times 48$  bits =  
 $2^{15} \times (1.5 \times 32)$  bits =  $1.5 \times 2^{20}$  bits = 1.5 Mbits
  - data bits in cache = 128 KB  $\times$  8 = 1 Mbits
  - total cache size/actual cache data = 1.5

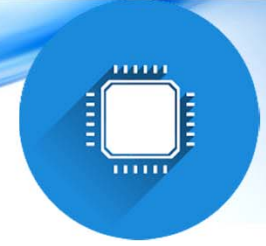




## مثال از نگاشت مستقیم

هر کلمه ۴ بایت است

- *How many total bits are required for a direct-mapped cache with 128 KB of data and 4-word block size, assuming a 32-bit address?*
- Cache size = 128 KB =  $2^{17}$  bytes =  $2^{15}$  words =  $2^{13}$  blocks
- Cache entry size = block data bits + tag bits + valid bit  
=  $128 + (32 - 13 - 2 - 2) + 1 = 144$  bits
- Therefore, cache size =  $2^{13} \times 144$  bits =  
 $2^{13} \times (1.25 \times 128)$  bits =  $1.25 \times 2^{20}$  bits = 1.25 Mbits
  - data bits in cache = 128 KB  $\times$  8 = 1 Mbits
  - total cache size/actual cache data = 1.25



## نگاشت مجموعه انجمنی (اشتراکی)

